

Equilibria Extremal Optimization

Rodica Ioana Lung

Facultatea de Științe Economice și Gestiunea Afacerilor

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- 1 Game theory
- 2 Nash Extremal Optimization
- 3 Application...

Solution concepts

Problem description:

- ▶ Game \rightarrow players, strategies, payoffs;
- ▶ Multiobjective optimization \rightarrow optimize multiple, conflicting objectives;

Solution concepts:

- ▶ Nash equilibrium (NE): no player can improve its payoff by unilateral deviation
- ▶ Berge equilibrium (BE): players maximize each others payoffs
- ▶ Pareto optimal solution/equilibrium: no objective can be increased without decreasing another objective

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Nash ascendancy relation

Compare two strategy profiles:

- ▶ compute $k(s, q) = \text{card}\{i \in N \mid u_i(q_i, s_{-i}) > u_i(s), q_i \neq s_i\}$
- ▶ s Nash ascends q if $k(s, q) < k(q, s)$
- ▶ s non-dominated with respect to the Nash ascendancy relation: $\nexists q \in S$ such that q Nash ascends s .
- ▶ Nash non-dominated = NE.

Probabilistic Nash ascendancy relation

Reduce the number of payoff function evaluations:

- ▶ only a percent p of players are tested:
- ▶ $I_p = \{i_1, i_2, \dots, i_{n_p}\} \subset N$
- ▶ then $k_p(s, q, I_p) = \text{card}\{i \in I_p \mid u_i(s) < u_i(q_i, s_{-i}), s_i \neq q_i\}$.
- ▶ $s \in S$ p -Nash ascends $q \in S$ with respect to I_p if $k_p(s, q, I_p) < k_p(q, s, I_p)$.
- ▶ $s \in S$ p -non-dominated $\nexists q \in S, \nexists I_p \subset N$ such that q p -ascends s with respect to I_p .
- ▶ p -non-dominated = Nash equilibria

Berge generative relation

Definition (Berge-Zhukovskii)

A strategy profile $s^* \in S$ is a Berge-Zhukovskii equilibrium if the inequality

$$u_i(s^*) \geq u_i(s_i^*, s_{-i})$$

holds for each player $i = 1, \dots, n$, and all $s_{-i} \in S_{-i}$.

Quality measures:

$$b(s, q) = \text{card}\{i \in N, u_i(s) < u_i(s_i, q_{-i}), s_{-i} \neq q_{-i}\},$$

$$b_\epsilon(s, q) = \text{card}\{i \in N, u_i(s) < u_i(s_i, q_{-i}) + \epsilon, s_{-i} \neq q_{-i}\},$$

The same mechanism \rightarrow generative relation

Evolutionary detection - Nash equilibria

Game:

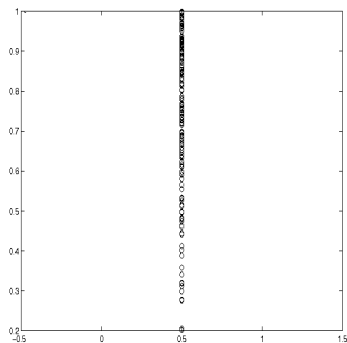
▶ Payoffs:

$$u_1(y_1, y_2) = y_1$$

$$u_2(y_1, y_2) = (0.5 - y_1)y_2$$

- ▶ $y_1 \in [0, 0.5]$, $y_2 \in [0, 1]$.
- ▶ NE: $(0.5, \lambda)$, $\lambda \in [0, 1]$.

NSGA-II results



Evolutionary detection - Berge Zhukovskii equilibria

Example

$$u_1(s_1, s_2) = -s_1^2 - s_1 + s_2,$$

$$u_2(s_1, s_2) = 2s_1^2 + 3s_1 - s_2^2 - 3s_2,$$

$$s_i \in [-2, 1], i = 1, 2.$$

Berge-Zhukovskii equilibrium: $(1, 1)$;

Payoffs $(-1, 1)$.

$\rightarrow \epsilon \in \{0, 0.1, 0.2, 0.5, 0.9\}$.

Cournot oligopoly

- ▶ $q_i, i = 1, \dots, N$ - quantities/ N companies
- ▶ market clearing price $P(Q) = a - Q$, where Q is the aggregate quantity on the market.
- ▶ total cost for the company i of producing quantity q_i :
 $C(q_i) = cq_i$.
- ▶ payoff for the company i is its profit:

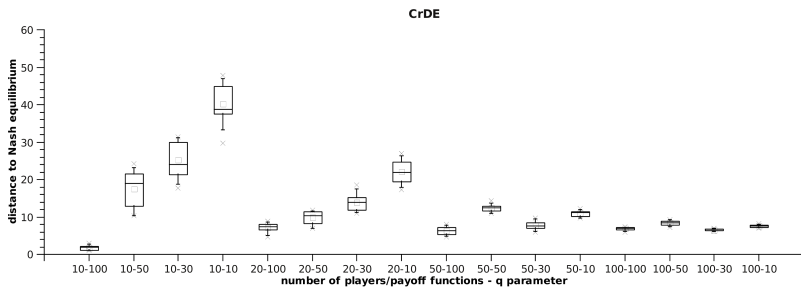
$$u_i(q_1, q_2, \dots, q_N) = q_i P(Q) - C(q_i)$$

- ▶ if $Q = \sum_{i=1}^N q_i$, the Cournot oligopoly has one Nash equilibrium

$$q_i = \frac{a - c}{N + 1}, \forall i \in \{1, \dots, N\}.$$

Multi-player games - probabilistic Nash ascendancy

Differential evolution



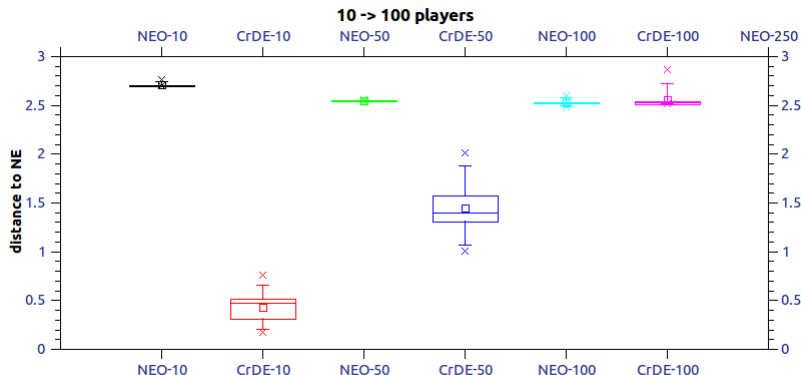
Nash Extremal Optimization

Algorithm 1 Nash Extremal Optimization procedure

- 1: Initialize configuration $s = (s_1, \dots, s_n)$ at will; set $s_{best} := s$;
 - 2: **repeat**
 - 3: For the 'current' configuration s evaluate u_i for each player i ;
 - 4: find j satisfying $u_j \leq u_i$ for all $i, i \neq j$, i.e., j has the "worst payoff";
 - 5: change s_j randomly;
 - 6: **if** (s **Nash ascends** s_{best}) **then**
 - 7: set $s_{best} := s$;
 - 8: **end if**
 - 9: **until** a termination condition is fulfilled;
 - 10: Return s_{best} .
-

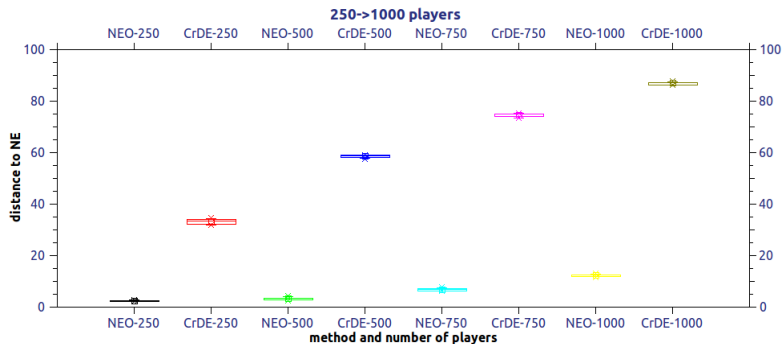
Multi-player games

Nash Extremal Optimization - Cournot

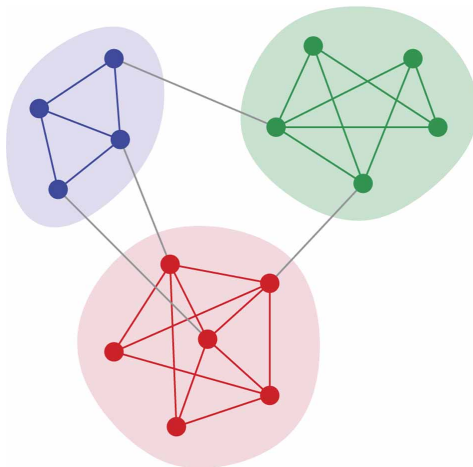


Multi-player games

Nash Extremal Optimization - Cournot

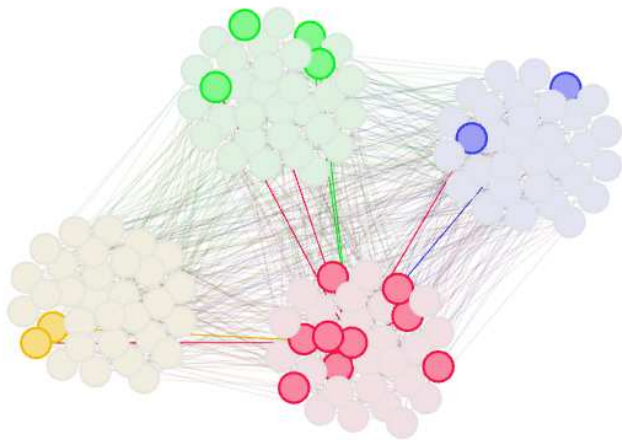


Community structure detection in social networks



M. E. J. Newman, Communities, modules and large-scale structure in networks, Nature Physics 8, 2531 (2012)
doi:10.1038/nphys2162

Community structure - definitions



The game

Game $\Gamma = (N, S, U)$:

- ▶ $N = \{1, \dots, n\}$, the set of players: network nodes;
- ▶ $S = S_1 \times S_2 \times \dots \times S_n$, the set of strategy profiles of the game; $s \in S$, $s = (c_1, \dots, c_n)$
- ▶ $U = \{u_i\}_{i=\overline{1,n}}$, the payoff functions; $u_i : S \rightarrow \mathbb{R}$ represents the payoff of player i , $i = \overline{1, n}$.

$$u_i(s) = f(C_i) - f(C_i \setminus \{i\}),$$

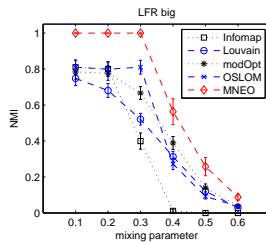
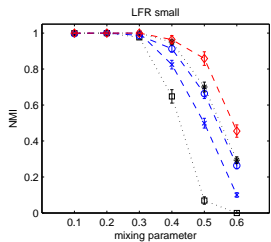
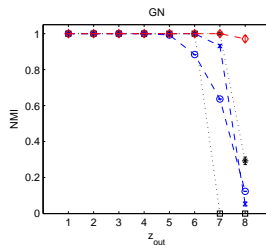
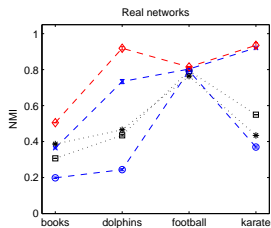
$$f(C) = \frac{k_{in}(C)}{(k_{in}(C) + k_{out}(C))^\alpha},$$

- ▶ $k_{in}(C)$ - double of the number of internal links in the community;
- ▶ $k_{out}(C)$ - number of external links of the community;
- ▶ α - a parameter controlling the size of the community.

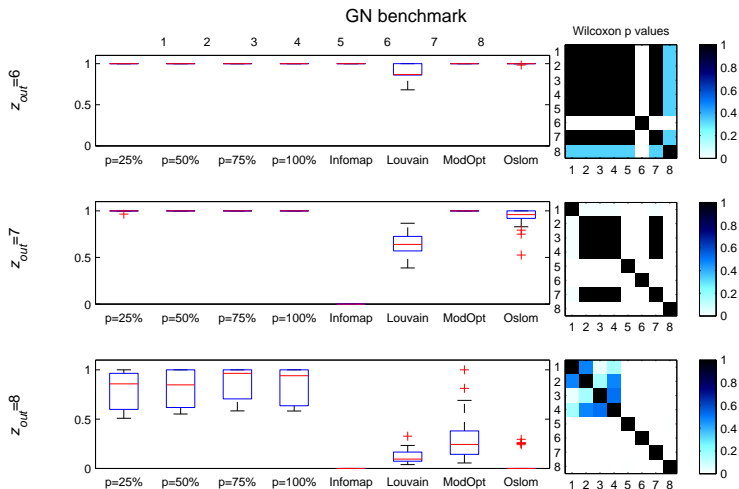
Mixed Network Extremal optimization

Randomly initialize all individuals in P and A ;
 Evaluate all individuals in P and A ;
for $NrGen = 0$ to $MaxGen$ **do**
 Set $m_{ngen} = \max \left\{ 1, \left[\frac{1}{10} \cdot N \cdot (N - 2)^{-\frac{NrGen}{MaxGen}} \right] \right\}$
 Run a **NEO** iteration for each pair (P_i, A_i) ;
 if Λ iterations were performed on the original network **then**
 Mix network with probability ρ ;
 Randomly re-initialize population A ;
 Evaluate all individuals in P and A ;
 end if
 if Network is mixed and λ iterations were performed **then**
 Restore network to original structure;
 Randomly re-initialize population A ;
 Evaluate all individuals in P and A ;
 end if
end for
Output: individual from A having the best modularity;

Results - MNEO

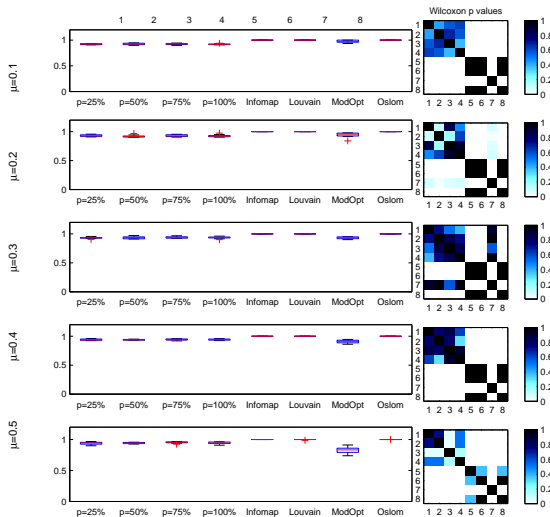


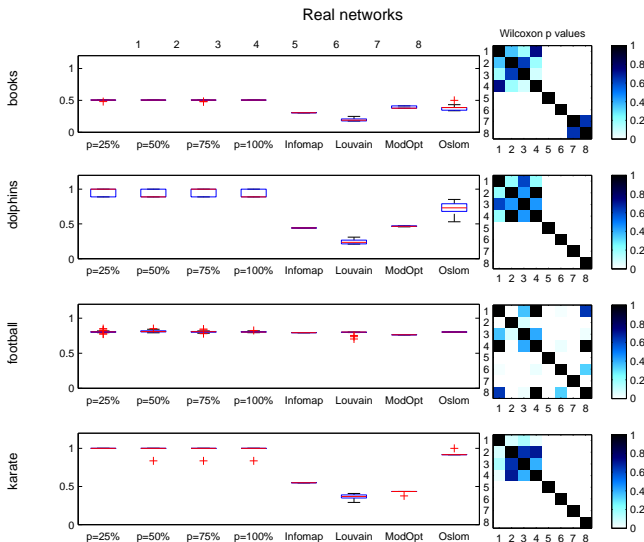
pMNEO



pMNEO

LFR benchmark, 1000 nodes, small



p MNEO

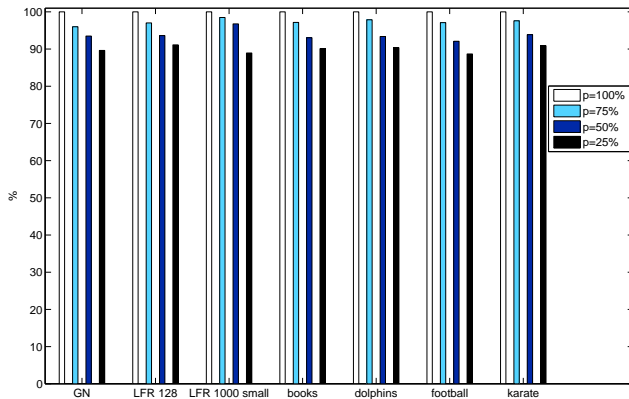
p MNEO

Figure: p MNEO duration in percents relative to $p = 100\%$

ϵ -Berge equilibria & the community detection problem

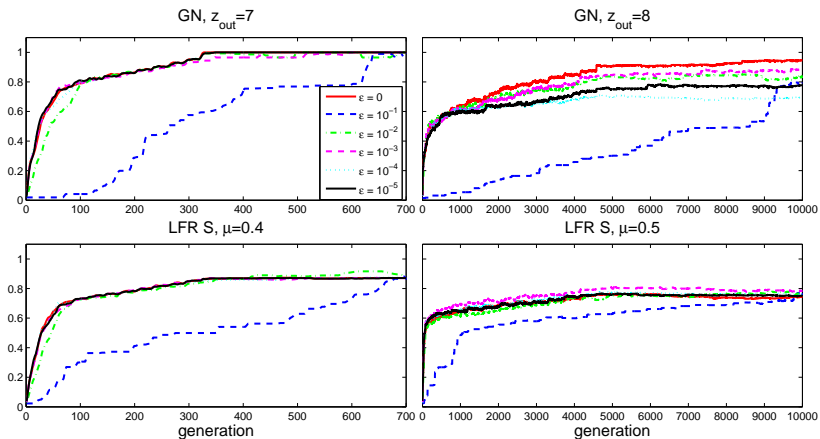


Figure: Evolution of NMI for different values of ϵ

ϵ -Berge equilibria & the community detection problem

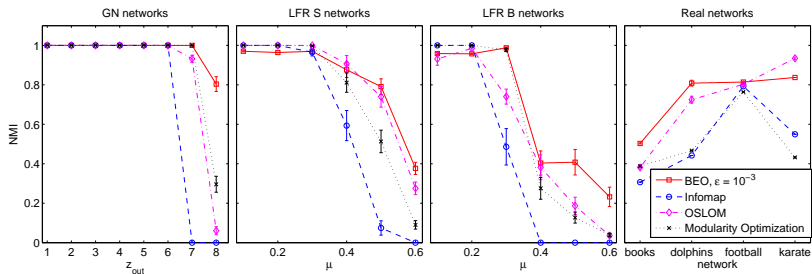


Figure: Berge equilibria - comparisons with other methods

Conclusions:

- ▶ (some) equilibria can be characterized by the use of "generative relations";
- ▶ generative relations can be embedded in evolutionary algorithms to compute corresponding equilibria;
- ▶ Extremal optimization is efficient for multi-player games;
- ▶ The community detection problem can be approached as a game;
- ▶ EO results are promising with both NE and BZ!

The work presented here was the result of the collaboration with:

- ▶ Noémi Gaskó
- ▶ Mihai Alexandru Suciú
- ▶ Tudor Dan Mihoc
- ▶ prof. D. Dumitrescu

Thank you for your attention!

Questions?

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